

Learning Goals

- Define variance and standard deviation
- State Chebyshev inequality and Chernoff inequality
- Compare the conditions and strengths of Markov, Chebyshev and Chernoff inequalities
- Understand the main idea and steps in the proofs of these bounds
- Intuition for the bounds given by the simplified forms of the Chernoff bound

Chebyshev Inequality

Definition

The *variance* of a random variable X is

$\text{Var}[X] := \mathbf{E}[(X - \mathbf{E}[X])^2] = \mathbf{E}[X^2] - (\mathbf{E}[X])^2$. Its square root, $\sqrt{\text{Var}[X]}$, is the *standard deviation* of X , and is often denoted as σ .

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Proof.

Apply Markov inequality to the random variable $(X - \mathbf{E}[X])^2$:

$$\Pr[|X - \mathbf{E}[X]| \geq \alpha\sigma] = \Pr[(X - \mathbf{E}[X])^2 \geq \alpha^2 \text{Var}[X]] \leq \frac{1}{\alpha^2}.$$



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A distribution where $|X - \mathbf{E}[X]|$ takes two values: 0 and $\alpha\sigma$
 $\Rightarrow X$ takes three values: $\mathbf{E}[X]$, $\mathbf{E}[X] + \alpha\sigma$ and $\mathbf{E}[X] - \alpha\sigma$.

Useful Facts for Independent Random Variables

Lemma

If X and Y are independent random variables, then $\mathbf{E}[XY] = \mathbf{E}[X] \cdot \mathbf{E}[Y]$, and $\text{Var}[X + Y] = \text{Var}[X] + \text{Var}[Y]$.

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$$\mathbf{E}[XY] = \sum_{x,y} (xy) \mathbf{Pr}[X = x, Y = y]$$

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Without independence, $\text{Var}[X + Y]$ in general is not equal to $\text{Var}[X] + \text{Var}[Y]$. □

Application of Chebyshev Inequality: Weak Law of Large Numbers

Theorem

Let X_1, X_2, \dots be *independently, identically distributed* (i.i.d.) random variables, and each has finite variance. For each $n \geq 1$, let \bar{X}_n be $\frac{1}{n} \sum_{i=1}^n X_i$. Then for any $\delta > 0$, $\lim_{n \rightarrow \infty} \Pr[|\bar{X}_n - \mathbf{E}[\bar{X}_n]| > \delta] = 0$.

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By Chebyshev inequality, $\Pr[|\bar{X}_n - \mathbf{E}[\bar{X}_n]| > \delta] \leq \frac{\text{Var}[X_1]}{n\delta^2}$.

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The right hand side goes to 0 as n goes to infinity. □

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 - Chebyshev \rightarrow Markov \rightarrow Kolmogorov
 - Bernstein and Chernoff exploited the idea by looking at $f(x) = e^{\lambda x}$.

Chernoff Bound: I.I.D. Case

Let X_1, \dots, X_n be i.i.d. Bernoulli variables, such that $\Pr[X_i = 1] = p$ and $\Pr[X_i = 0] = q := 1 - p$ for each i . Define $X = \sum_{i=1}^n X_i$.

Theorem (Chernoff Bound)

For any $t > 0$,

$$\Pr[X > (p + t)n] \leq \exp \left\{ \left(-(p + t) \ln \frac{p + t}{p} - (q - t) \ln \frac{q - t}{q} \right) n \right\}.$$

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The same proof yields the same bound for $\Pr[X \leq (p - t)n]$.

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Useful Forms of Chernoff Bound

Corollary

Let X_1, \dots, X_n be independently distributed on $[0, 1]$ and $X = \sum_i X_i$.

- For all $t > 0$,

$$\Pr[X > \mathbf{E}[X] + t], \Pr[X < \mathbf{E}[X] - t] \leq e^{-2t^2/n};$$

- For any $\epsilon < 1$,

$$\Pr[X > (1 + \epsilon) \mathbf{E}[X]] \leq \left(\frac{e^\epsilon}{(1 + \epsilon)^{1+\epsilon}} \right)^{\mathbf{E}[X]} \leq \exp \left(-\frac{\epsilon^2}{3} \mathbf{E}[X] \right);$$

$$\Pr[X < (1 - \epsilon) \mathbf{E}[X]] \leq \left(\frac{e^{-\epsilon}}{(1 - \epsilon)^{1-\epsilon}} \right)^{\mathbf{E}[X]} \leq \exp \left(-\frac{\epsilon^2}{2} \mathbf{E}[X] \right).$$

Useful Forms of Chernoff Bound (Cont.)

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- For any $\epsilon > 1$,

$$\Pr [X > (1 + \epsilon) \mathbf{E} [X]] \leq \left(\frac{e^\epsilon}{(1 + \epsilon)^{1+\epsilon}} \right)^{\mathbf{E}[X]} \leq \exp \left(-\frac{\epsilon}{3} \mathbf{E} [X] \right);$$

- If $t > 2e \mathbf{E}[X]$, then

$$\Pr [X > t] \leq 2^{-t}.$$

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- Let $f(t)$ be $(p+t) \ln \frac{p+t}{p} + (q-t) \ln \frac{q-t}{q}$. Show $f(t) \geq 2t^2$ by showing $f(0) = f'(0) = 0$ and $f''(t) \geq 4$ for all $0 \leq t \leq q$ followed by Taylor's theorem with remainder.

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- Let $g(x)$ be $f(px)$, then $g'(0) = pf'(px)$, and so $g(0) = g'(0) = 0$. Show $g'(1) > p \ln 2 > \frac{2}{3}p$. Deduce that for $x \in (0, 1)$, $g(x) \geq px^2/3$.

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- Let $g(x)$ be $f(px)$, then $g'(0) = pf'(px)$, and so $g(0) = g'(0) = 0$. Show $g'(1) > p \ln 2 > \frac{2}{3}p$. Deduce that for $x \in (0, 1)$, $g(x) \geq px^2/3$.
- Set $h(x) := g(-x)$. Then $h'(x) = -g'(-x)$, and $h(0) = h'(0) = 0$. Show then $h''(x) \leq p$ for $x \in (0, 1)$. Deduce that $h(x) \geq px^2/2$.

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See assigned reading for more details. Or take them as exercises.

